

From FDWs to Sharding

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Pushdown

- SELECT * FROM ft1 WHERE x = 1 → Yes!
- SELECT * FROM ft1 ORDER BY y → No.
- SELECT * FROM ft1, ft2 WHERE ft1.x = ft2.x \rightarrow No.
- SELECT count(*) FROM ft1 → No!!
- UPDATE ft1 SET x = 1 WHERE $y = 2 \rightarrow No!!$



Partitioning: No!

Join

- → Append
 - → Scan server s1, table x1
 - → Scan server s2, table x2
 - → Scan server s3, table x3
- → Append
 - → Scan server s1, table y1
 - → Scan server s2, table y2
 - → Scan server s3, table y3



Partitioning: Yes!

Append

- → Join
 - → Scan server s1, table x1
 - → Scan server s1, table y1
- → Join
 - → Scan server s2, table x2
 - → Scan server s2, table y2
- → Join
 - → Scan server s3, table x3
 - → Scan server s3, table y3



Partitioning: Aggregates

AggregateCombiner

- → PartialAggregate
 - → Scan server s1, table x1
- → PartialAggregate
 - → Scan server s2, table x2
- → PartialAggregate
 - → Scan server s3, table x3



Partitioning: Replication

SELECT * FROM sharded_table JOIN codes_table ON sharded_table.codes_table_id = codes_table.id

Maybe there's a copy of codes_table on every node!



MVCC

Atomic Commit

 Single update touches multiple shards, either all updates commit or all updates roll back.

Atomic Visibility

 Single update touches multiple shards, no one can see an intermediate state where some but not all updates have committed.

